

## Arc Consistency and AC-3

$X_i$  is **arc consistent** (2-consistent) w.r.t.  $X_j$  if for all values  $v \in D_i$  there exists  $w \in D_j$  s.t.  $(v,w) \in C(X_i, X_j)$ .

*If the CSP is arc consistent then a solution exist? Not necessarily.*

If AC-3 terminates with an empty domain  $\Rightarrow$  **no solution exists**

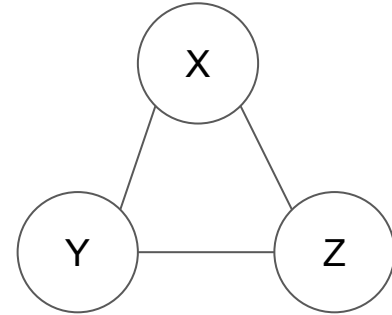
If AC-3 terminates with non empty domains  $\Rightarrow$  **a solution might exist**

# Path Consistency

*3 Map coloring problem*

$D(X) = D(Y) = D(Z) = \{\text{red}, \text{blue}\}$

$C(X,Y) = C(X,Z) = C(Y,Z) = \{(\text{red},\text{blue}),(\text{blue},\text{red})\}$



**Arc consistent but no solution exists!  $\Rightarrow$  Not path consistent**

$\{X_i, X_j\}$  is **path consistent** w.r.t.  $X_k$  if for all values  $v \in D_i$ ,  $w \in D_j$ , and  $(v,w) \in C(X_i, X_j)$  there exists  $t \in D_k$  s.t.  $(v,t) \in C(X_i, X_k)$  and  $(w,t) \in C(X_j, X_k)$ .

**A path consistent CSP can still have no solution!** *4 Map coloring problem*

# K-consistency

A CSP is **K-consistent** if for every subset of  $K-1$  variables, consistent assignment to those variables and for every  $K$ -th variable  $Y$ , there exists a consistent assignment for  $Y$ .

- 1-consistent = node consistent
- 2-consistent = arc consistent
- 3-consistent + all constraints are binary = path consistent

A CSP is **strongly K-consistent** if it is  $J$ -consistent for all  $1 \leq J \leq K$ .

**Let  $N$  be the number of variables of a CSP. If the CSP is strongly  $N$ -consistent then a solution exists.**